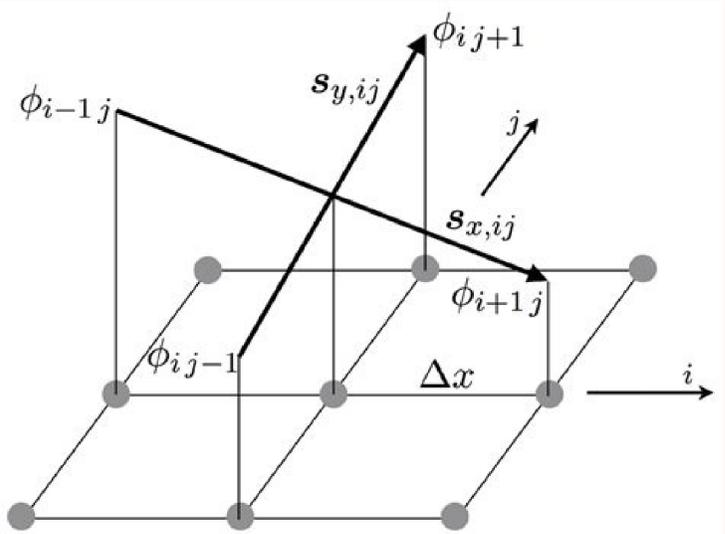
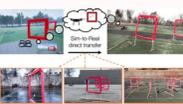
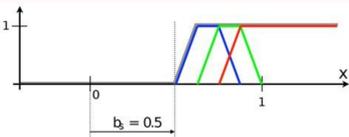
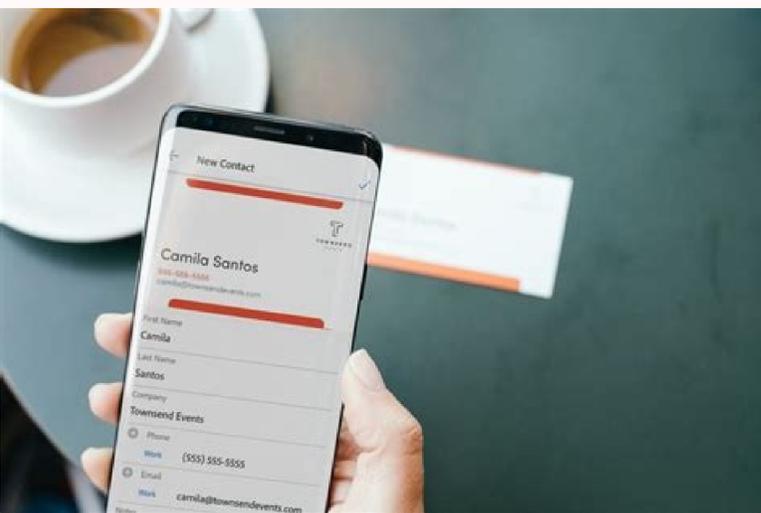


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In particular, number restrictions are not used. In this way, we also reprove in an elegant way Danecki's difficult result that satisfiability for PDL with intersection is in 2EXP. • A domino must be placed exactly over a ruled square. Moreover, only soundness and completeness of the tableau algorithm have to be proved; termination and a practical procedure to ensure termination, called a blocking test, also follows from the prerequisites of the framework. Note: Here EXP means ExpTime, and 2EXP means ExpTime. The proof uses a reduction from the word problem for alternating, exponential time bounded Turing machines. E.Z.: In short: the converse μ -calculus is ExpTime-complete. Information and Computation, 199(1-2):132-171, 2005. This strengthens the results by Grädel, Kolaitis, and Vardi [15], who showed that the satisfiability problem for the first-order two-variable logic L2 is NExpTime-complete and by Grädel, Otto, and Rosen [16], who proved the decidability of C2. E.Z.: In short, ICPDL and hence ALCI ($\cap, \cup, \sigma, \text{id}$) is decidable. One relevant and powerful variation is ICPDL, the extension of PDL with intersection and converse. Eliminating "converse" from converse PDL [pdf] @article{DeGiacomoLenzerini1996a, author = {Giuseppe De Giacomo}, title = {Eliminating "converse" from converse PDL}, journal = {Journal of Logic, Language, and Information}, volume = {5}, number = {2}, year = {1996}, pages = {193-208}, doi = {10.1007/BF00173700}, ee = { } In this paper we show that it is possible to eliminate the "converse" operator from the propositional dynamic logic CPDL (Converse PDL), without compromising the soundness and completeness of inference for it. On the other hand, with the rather weak form of number restrictions available in implemented systems, the number of role successors of an individual can only be restricted by a fixed non-negative integer. Hence, so is concept satisfiability in μ ALCI. LNCS, vol. 4548, pp. 1 allow only nonperiodic tilings of the plane. We refer to [DL Handbook] for usual background knowledge on description logics (DLs). If the number of modal parameters is bounded, full Boolean Modal Logic becomes ExpTime-complete. The expressive power of this constructor is increased even further by introducing explicit quantifiers for the numerical variables. It is worth noticing that by a recent result of E. A recent paper by Schild showed that this issue can be profitably addressed by relying on a correspondence between DLs and propositional dynamic logics (PDL). It would be desirable to combine different maximal logics into one super-logic, but then inference may turn out to be undecidable. A tableaux decision procedure for SHOIQ [pdf] @inproceedings{HorrSattlerSHOIQ2005, author = {Ian Horrocks and Ulrike Sattler}, title = {A tableaux decision procedure for Σ matcal{SHOIQ}}, pages = {448-453}, editor = {Leslie Pack Kaelbling and Alessandro Saffioti}, booktitle = {Proc. of the 19th Int. Joint Conf. on Artificial Intelligence (IJCAI-2005)}, publisher = {Professional Book Center}, year = {2005} } This paper presents a tableaux decision procedure for SHOIQ, the DL underlying OWL-DL, vol. 106 in CEUR, pp. Finally, we show how, combining the technique of 'filtration' with a notion of 'copying' worlds - in view of the "more than n" interpretation, one cannot simply collapse worlds - for some modal logics (GrK, GrT, ...), the finite model property (and also decidability) is obtained. Theoretical Computer Science, Special Issue: 9th Int. Shapiro, editors, Proc. Technical report: Is your RBox safe? The original motivation for this work was to develop a description logic suitable for qualitative spatial reasoning problems. Technical report: [pdf] A PSpace-algorithm for deciding ALCNI R4-satisfiability [ps.zip] @techreport{HorrocksSattlerTobies1998, author = {Ian Horrocks and Ulrike Sattler and Stephan Tobies}, title = {A PSpace-algorithm for deciding Σ matcal{ALCNI} (R^+)-satisfiability}, institution = {LuFg Theoretical Computer Science, RWTH Aachen}, address = {Germany}, type = {LTCS-Report}, number = {98-08}, year = {1998} } ALCNI R+ - ALCN augmented with transitive and inverse roles - is an expressive Description Logic which is especially well-suited for the representation of complex, aggregated objects. This fact allows one to build inference procedures for CPDL, by encoding CPDL formulae into PDL, and then running an inference procedure for PDL. Second, we allow to express inclusion statements between general concepts, and terminological cycles as a particular case. The second is our older idea to formalize the notion of 'object' in the same modal spirit. Secondly, we will see that although these two logics share most of their constructs, they lead to quite different TSS, which demonstrates how the capabilities of our framework can be used to capture different language properties. It is interesting that C12 is in NExpTime in spite of the fact that there are sentences whose minimal (and only) models are of doubly exponential size, of the 2002 Int. on the Principles of Knowledge Representation and Reasoning, KR 2006, Lake District of the UK, June 2-5, 2006, 122-131, 2001. In contrast, intuitive automata algorithms usually require exponential time also in the best case, and thus an efficient implementation requires considerable modifications. 253-263, Plenum Press, New York, 1987. The intersection, cf. In Harald Ganzinger, David McAllester, and Andrei Voronkov, editors, Proc. We see to assemble the copies (also called dominoes) on an infinite plane, ruled into squares the size of one domino, according to the following rules: • No domino may be rotated or reflected. We propose to further extend such DLs with key constraints that allow the expression of statements like "US citizens are uniquely identified by their social security number". We present an algorithm that decides satisfiability of the DL ALC extended with transitive and inverse roles and functional restrictions with respect to general concept inclusion axioms and role hierarchies; early experiments indicate that this algorithm is well-suited for implementation. In Proc. In L.C. Aiello, Jon Doyle, and Stuart C. Morgan Kaufmann, 2003. of the 5th Symposium on Computation Theory, Zaborów, Poland, December 3-8, 1984, LNCS, vol. 208, pp. Information and Computation, 81(3):249-264, 1989. In Artificial Intelligence, AAAI'94, Seattle, Washington, July 31 - August 4, 1994, pp. Workshop on Description Logics, DL 2006, Windermere, Lake District, UK, May 30 - June 1, 2006 vol. 189 in CEUR, pp. We provide applications of the decision procedure to regular expressions, Ianov schemes, and classical systems of modal logic. Joint Conf. Journal of Artificial Intelligence Research, vol. 6, pp. Journal of Artificial Intelligence Research, 23:667-726, 2004. Finally, we consider negated role assertions in ABoxes and qualified number restrictions. 540-551, Springer, 2006. Technical report: [pdf] Decidability of expressive Description Logics with role compositions (extended abstract) [ps.zip] In Proc. The complexity of reasoning with cardinality restrictions and nominals in expressive Description Logics [pdf] @article{Tobies2000, author = {Stephan Tobies}, title = {The complexity of reasoning with cardinality restrictions and nominals in expressive Description Logics}, journal = {Journal of Artificial Intelligence Research}, volume = {12}, year = {2000}, pages = {199-217}, doi = {10.1613/jair.705}, ee = { } We study the complexity of the combination of the Description Logics ALCQ and ALCQI with a terminological formalism based on cardinality restrictions on concepts. Such an extension of Km is called a regular grammar logic. Vardi, On the decision problem for two-variable first-order logic. In L.C. Aiello, Jon Doyle, and Stuart C. Shapiro, editors, Proc. We analyze from a theoretical point of view a TKRS whose capabilities go beyond the ones of presently available TKRSs. The new features studied, often required in practical applications, can be summarized in three main points. The contribution of the paper is two-fold: we extend the family of ExpTime logics, and we present a technique to reason in the presence of nominals. Decidable reasoning in terminological knowledge representation systems [pdf] @article{Buchheit et al 1993a, author = {Martin Buchheit and Francesco M. Decidability of class-based knowledge representation formalisms [pdf] [ps.zip] PhD Thesis, Dipartimento di Informatica e Sistemistica, Università di Roma "La Sapienza", 1995. By discovering that features correspond to deterministic programs in dynamic logic, we show that adding them to TSL preserves decidability, although violates its finite model property. In Rohit Parikh, editor, Proc. An abridged version appeared in Proc. We show that this approach leads to an optimal solution for ALCQI, as ALCQI with cardinality restrictions has the same complexity as C2 (NExpTime-complete). Shapiro), booktitle = {Proc. of the 5th Int. Conf. on Principles of Knowledge Representation and Reasoning (KR'96)}, publisher = {Morgan Kaufmann}, year = {1996} } A Description Logic (DL) system is characterized by four fundamental aspects: the set of constructs used in concept and role expressions, the kind of assertions allowed in the TBox (assertions on concepts) and the ABox (assertions on individuals), and the inference mechanisms for reasoning on both the TBox and the ABox. Our result has applications in information logic, description logic, and epistemic logic. (a) Proof that the six polygonal tiles shown on Fig. Hence, concept satisfiability in μ ALCQI is ExpTime-complete (even with general TBoxes, which are internalizable in presence of μ). Additionally, we describe an algorithm for deciding the coherence of inverse-free TSL-concepts with features. 168-177, 2004. Technical report: [ps.zip] LTCS-Report 03-04, Chair for Automata Theory, Institute for Theoretical Computer Science, Dresden University of Technology, Germany, 2003. Nebel, editor, Proc. One of the connections made between the models, that involving language models, is quite counterintuitive. From the standpoint of DLs, we derive decidability and complexity results for some of the most expressive logics appeared in the literature, and from the standpoint of PDLs, we derive a general methodology for the representation of several forms of program determinism and for the specification of partial computations. (574 pages) ISBN 978-0-521-78176-3 Handbook of Modal Logic Modal Logic Computational Complexity Addison-Wesley Publishing Company, 1994. 1200, Springer, Berlin, 1997. 57-67, AAAI Press, 2006. To make it even more useful for applications, many extensions of PDL have been considered in the literature. of the 12th IEEE Symposium on Logic in Computer Science, LICS'97, Warsaw, Poland, June 29 - July 2, 1997, pp. Expressive role constructors are important in many applications, but can be computationally problematical. Note: This Description Logic is now known as RIQ. In comparison with PDL with strong loop predicate (Danecki 1984), this is more powerful and interesting example of a logic which is decidable in contrast to its deterministic case and despite the lack of finite and even tree model properties. On the one hand, the roles that may occur in number restrictions are usually of a very restricted type, namely atomic roles or complex roles built using either intersection or inversion. While PDL consists of assertions about weakest preconditions, the converse construct enable us to make assertions about strongest postconditions. Surprisingly, this extension is still undecidable. Vardi), title = {Reasoning about the past with two-way automata}, pages = {628-641}, editor = {Kim G. In Dimitar G. SHOQ (D) is an expressive description logic equipped with named individuals and concrete datatypes which has almost exactly the same expressive power as the latest web ontology languages (e.g., OIL and DAML+OIL), of the 14th Annual Symposium on Theoretical Aspects of Computer Science, LNCS, vol. A lower complexity bound for Propositional Dynamic Logic with intersection [pdf] [ps.zip] @inproceedings{LangeEXPSPACE, author = {Martin Lange}, title = {A lower complexity bound for Propositional Dynamic Logic with intersection}, pages = {133-147}, editor = {Renate A. Instances of realistic size for PSpace-complete problems are now within reach for implemented systems. It considerably extends L2, plain first-order with only two variables, which is known to be decidable by a result of Mortimer. Augmenting concept languages by transitive closure of roles: An alternative to terminological cycles [pdf] @inproceedings{Baader1991, author = {Franz Baader}, title = {Augmenting concept languages by transitive closure of roles: An alternative to terminological cycles}, pages = {446-451}, editor = {John Mylopoulos and Raymond Reiter}, booktitle = {Proc. of the 12th Int. Joint Conf. on Artificial Intelligence (IJCAI'91)}, publisher = {Morgan Kaufmann}, year = {1991} } In Baader (1990, 1990a), we have considered different types of semantics for terminological cycles in the concept language TLQ which allows only conjunction of concepts and value-restrictions. Whereas tableau-based algorithms usually yield worst-case optimal algorithms in the case of PSpace-complete logics, it is often very hard to design optimal tableau-based algorithms for ExpTime-complete DLs. In contrast, the automata-based approach is usually well-suited to prove ExpTime upper-bounds, but its direct application will usually also yield an ExpTime-algorithm for a PSpace-complete logic since the (tree) automaton constructed for a given concept is usually exponentially large. For this reason, an approach combining both advantages is highly desirable. of the 18th Conference on Automated Deduction, CADE 2002, Copenhagen, Denmark, July 27-30, 2002, LNCS, vol. 2392, pp. Two-variable logic with counting is decidable [pdf] [ps.zip] @inproceedings{Gradel et al 1997b, author = {Erich Gradel and Martin Otto and Eric Rosen}, title = {Two-variable logic with counting is decidable}, pages = {306-317}, editor = {Glynn Winskel}, booktitle = {Proc. of the 12th Annual IEEE Symposium on Logic in Computer Science (LICS'97)}, publisher = {IEEE Computer Society Press}, year = {1997}, doi = {10.1109/LICS.1997.614957}, ee = { } We prove that the satisfiability problem for C2 is decidable. Finally, we prove that universal implications can be expressed within TSL. 704-709, Morgan Kaufmann, 1993. It is known that the satisfiability problem for the μ -calculus is ExpTime-complete. Research report: [ps.zip] Expressive number restrictions in Description Logics [pdf] [ps.zip] @article{BaaderSattler1999, author = {Franz Baader and Ulrike

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